

Finite model theory
 Problems 11
 Tuesday 27.11.2018

1. Let φ be a first-order formula with vocabulary τ ($\{\min, \max\} \subseteq \tau$) and \vec{x} and \vec{y} tuples of variables not appearing free in φ . Show that

$$[\text{TC}_{\vec{x}, \vec{y}} \varphi] \vec{\min} \vec{\max}$$

is equivalent to φ .

2. Let $\varphi(\vec{x}) \in \text{FO}(\text{posTC})$ be a formula. Construct a formula $\text{count}_\varphi(\vec{x}) \in \text{FO}(\text{posTC})$ such that for all ordered finite models \mathfrak{A} and \vec{a} :

$$\mathfrak{A} \models \text{count}_\varphi[\vec{a}/\vec{x}] \Leftrightarrow \vec{a} = |\{\vec{b} \mid \mathfrak{A} \models \varphi[\vec{b}/\vec{x}]\}|,$$

where $\vec{a} = (a_0, \dots, a_{k-1})$ is taken to encode the number

$$a_0 n^{k-1} + a_1 n^{k-2} + \dots + a_{k-1},$$

and where $\text{Dom}(\mathfrak{A}) = \{0, \dots, n-1\}$.

3. Let $\varphi = \varphi_1 \wedge \varphi_2$, where φ_i is of the form

$$[\text{TC}_{\vec{x}, \vec{y}} \psi_i] \vec{\min} \vec{\max}$$

and ψ_i is quantifier-free for $i = 1, 2$. Show that then φ is also equivalent to a formula of the form

$$[\text{TC}_{\vec{x}', \vec{y}'} \psi] \vec{\min} \vec{\max},$$

where ψ is quantifier-free.

4. Show that existential second-order logic Σ_1^1 is closed under \vee and $\forall x$ in the following sense: if $\varphi, \psi \in \Sigma_1^1$, then there is $\theta \in \Sigma_1^1$ such that for all \mathfrak{A} :

$$\mathfrak{A} \models \varphi \vee \psi \Leftrightarrow \mathfrak{A} \models \theta$$

and

$$\mathfrak{A} \models \forall x \varphi \Leftrightarrow \mathfrak{A} \models \theta.$$

5. Show that for ordered finite models: $\text{NPTIME} = \text{co-NPTIME}$ iff $\Sigma_1^1 \equiv \text{SO}$. Here, for each vocabulary τ , co-NPTIME contains the complements of NPTIME-classes:

$$\{K \subseteq O[\tau] \mid (O[\tau] \setminus K) \in \text{NPTIME}\}.$$