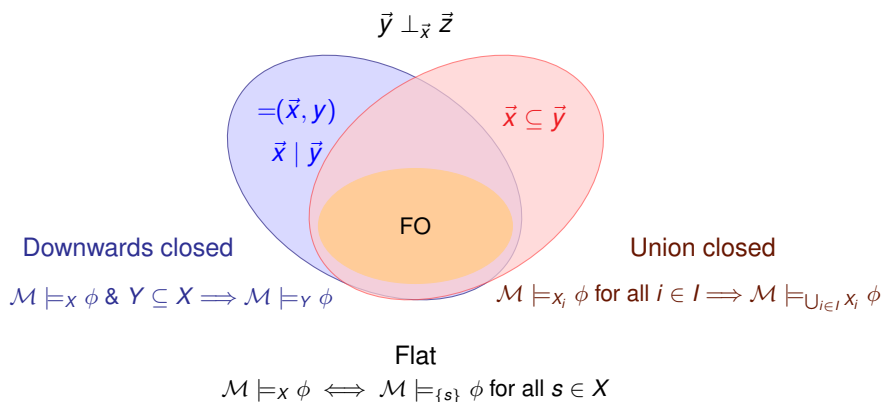


Dependence Logic

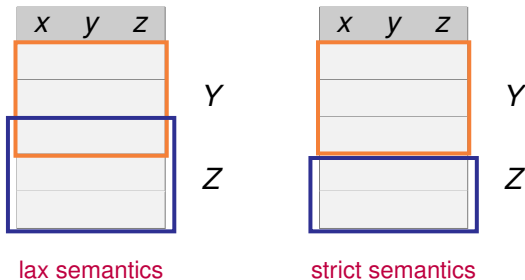
September 17 (Tuesday), 2019

Recap: Closure properties



Recap: Disjunction

- $\mathcal{M} \models_X \phi \vee \psi$ iff $\exists Y, Z \subseteq X$ s.t. $X = Y \cup Z$,
 $M \models_Y \phi$ and $M \models_Z \psi$.
- $\mathcal{M} \models_X^s \phi \vee \psi$ iff $\exists Y, Z \subseteq X$ s.t. $X = Y \cup Z$, $Y \cap Z = \emptyset$,
 $M \models_Y \phi$ and $M \models_Z \psi$.



Locality: For any two teams X, Y with $X \upharpoonright \text{Fv}(\phi) = Y \upharpoonright \text{Fv}(\phi)$,

$$\mathcal{M} \models_X \phi \iff \mathcal{M} \models_Y \phi.$$

Note: Locality fails for $\text{FO}(\subseteq)$ with strict semantics.

Existential quantifier

- $\mathcal{M} \models_X \exists x \phi$ iff $\mathcal{M} \models_{X(F/x)} \phi$ for some $F : X \rightarrow \wp^+(M)$, where $X(F/x) = \{s(a/x) \mid s \in X \text{ and } a \in F(s)\}$.
- $\mathcal{M} \models_X^s \exists x \phi$ iff $\mathcal{M} \models_{X(F_0/x)} \phi$ for some $F_0 : X \rightarrow M$, where $X(F_0/x) = \{s(F_0(s)/x) \mid s \in X\}$.

<i>u</i>	<i>v</i>	<i>x</i>
<i>a</i>	<i>b</i>	<i>a</i>
<i>b</i>	<i>c</i>	<i>b</i>
<i>b</i>	<i>d</i>	<i>c</i>
<i>c</i>	<i>a</i>	<i>d</i>
\vdots	\vdots	\vdots

M

lax semantics

<i>u</i>	<i>v</i>	<i>x</i>
<i>a</i>	<i>b</i>	<i>a</i>
<i>b</i>	<i>c</i>	<i>b</i>
<i>b</i>	<i>d</i>	<i>c</i>
<i>c</i>	<i>a</i>	<i>d</i>
\vdots	\vdots	\vdots

M

strict semantics

Locality: For any two teams X, Y with $X \upharpoonright \text{Fv}(\phi) = Y \upharpoonright \text{Fv}(\phi)$,

$$\mathcal{M} \models_X \phi \iff \mathcal{M} \models_Y \phi.$$

Note: Locality fails for $\text{FO}(\subseteq)$ with strict semantics.